## IP and Networking Basics

#### Scalable Infrastructure Workshop AfNOG 2011

#### 1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packet-switching
- 1964: Baran packetswitching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

**1972**:

- ARPAnet demonstrated publicly
- NCP (Network Control Protocol) first host-host protocol
- first e-mail program
- ARPAnet has 15 nodes

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1973: Metcalfe's PhD thesis proposes Ethernet
- 1974: Cerf and Kahn architecture for interconnecting networks
- Late 70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

- Cerf and Kahn's internetworking principles:
  - minimalism, autonomy no internal changes required to interconnect networks
  - best effort service model
  - stateless routers
  - decentralized control
- define today's Internet architecture

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: SMTP e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: FTP protocol defined
- 1988: TCP congestion control

- New national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

1990, 2000's: commercialisation, the Web, new apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)

early 1990s: Web

- hypertext [Bush 1945, Nelson 1960's]
- HTML, HTTP: Berners-Lee
- 1994: Mosaic, later Netscape
- late 1990's: commercialization of the Web

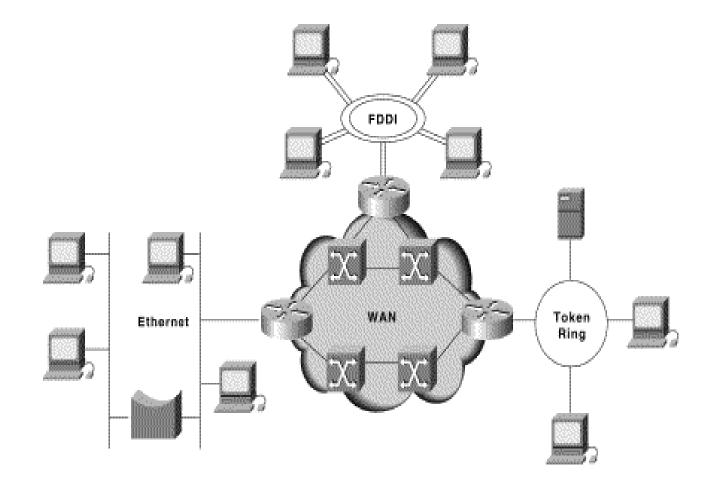
Late 1990's – 2000's:

- more killer apps: instant messaging, peer2peer file sharing (e.g., Napster)
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps
- now: 40-100 Gbps
  - youtube, social networking
  - depletion of Ipv4 address space

### The (capital "I") Internet

- The world-wide network of TCP/IP networks
- Different people or organisations own different parts
- Different parts use different technologies
- Interconnections between the parts
- Interconnections require agreements
  - sale/purchase of service
  - contracts
  - "peering" agreements
- No central control or management

## A small internetwork or (small "i") "internet"

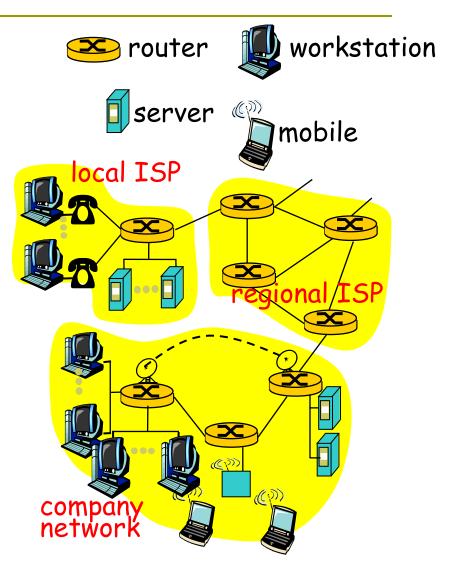


#### The principle of "Internetworking"

- We have lots of little networks
- Many different owners/operators
- Many different types
  - Ethernet, dedicated leased lines, dialup, optical, broadband, wireless, ...
- Each type has its own idea of low level addressing and protocols
- We want to connect them all together and provide a unified view of the whole lot (treat the collection of networks as a single large internetwork)

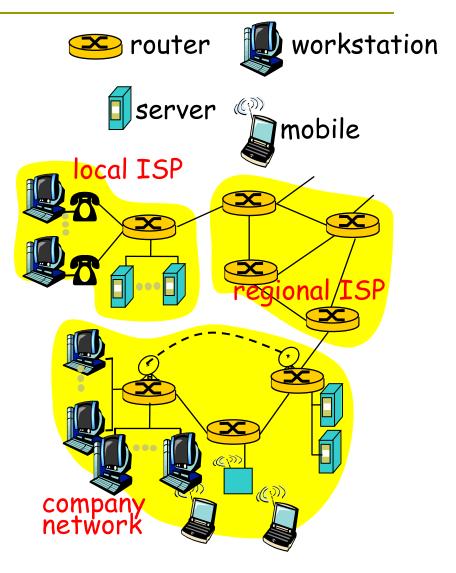
#### What is the Internet: "nuts and bolts" view

- millions of connected computing devices: hosts, end-systems
  - PC's workstations, servers
  - PDA's phones, toasters
  - running network apps
- communication links
  - fiber, copper, radio, satellite
- routers: forward packets (chunks) of data through network



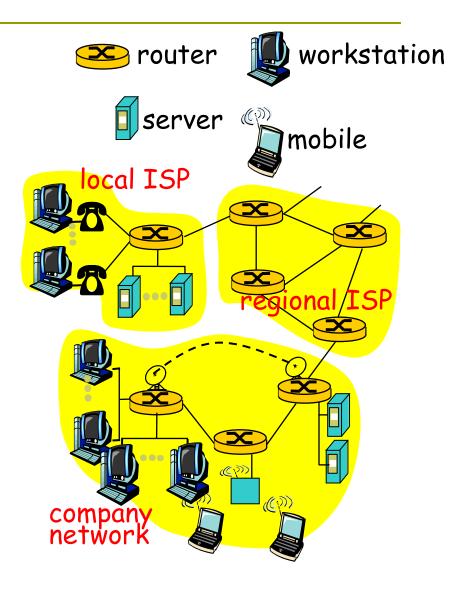
#### What is the Internet: "nuts and bolts" view

- protocols: control sending, receiving of messages
  - e.g., TCP, IP, HTTP, FTP, PPP
- Internet: "network of networks"
  - loosely hierarchical
  - public Internet versus private intranet
- Internet standards
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force



#### What is the Internet: a service view

- communication infrastructure enables distributed applications:
  - WWW, email, games, e-commerce, database, e-voting, more?
- communication services provided:
  - connectionless
  - connection-oriented



#### Connectionless Paradigm

#### There is no "connection" in IP

- Packets can be delivered out-of-order
- Each packet can take a different path to the destination
- No error detection or correction in payload
- No congestion control (beyond "drop")
- TCP mitigates these for connectionoriented applications
  - error correction is by retransmission
  - Packet drops as congestion signalling

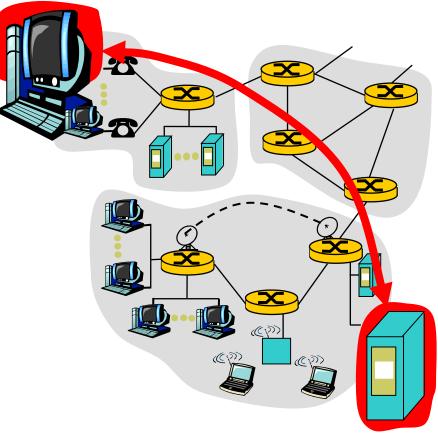
## OSI Stack & TCP/IP Architecture

#### Principles of the Internet

- Edge vs. core (end-systems vs. routers)
  - Dumb network
  - Intelligence at the end-systems
- Different communication paradigms
  - Connection oriented vs. connection less
  - Packet vs. circuit switching
- Layered System
- Network of collaborating networks

#### The network edge

- end systems (hosts):
  - run application programs
  - e.g., WWW, email
  - at "edge of network"
- client/server model:
  - client host requests, receives service from server
  - e.g., WWW client (browser)/server; email client/server
- peer-peer model:
  - host interaction symmetric e.g.: teleconferencing



# Network edge: connection-oriented service

- Goal: data transfer between end sys.
- handshaking: setup (prepare for) data transfer ahead of time
  - Hello, hello back human protocol
  - set up "state" in two communicating hosts
- TCP Transmission Control Protocol
  - Internet's connectionoriented service

- **TCP** service [RFC 793]
- reliable, in-order bytestream data transfer
  - Ioss: acknowledgements and retransmissions
- flow control:
  - sender won't overwhelm receiver
- congestion control:
  - senders "slow down sending rate" when network congested

# Network edge: connectionless service

Goal: data transfer between end systems

- UDP User Datagram Protocol [RFC 768]: Internet's connectionless service
  - unreliable data transfer
  - no flow control
  - no congestion control

#### Protocol "Layers"

- Networks are complex!
- many "pieces":
  - hosts
  - routers
  - links of various media
  - applications
  - protocols
  - hardware, software

#### Question:

- Is there any hope of organizing structure of network?
- Or at least in our discussion of networks?

The unifying effect of the network layer

- Define a protocol that works in the same way with any underlying network
- Call it the network layer (e.g. IP)
- IP routers operate at the network layer
- IP over anything
- Anything over IP

## Why layering?

#### Dealing with complex systems:

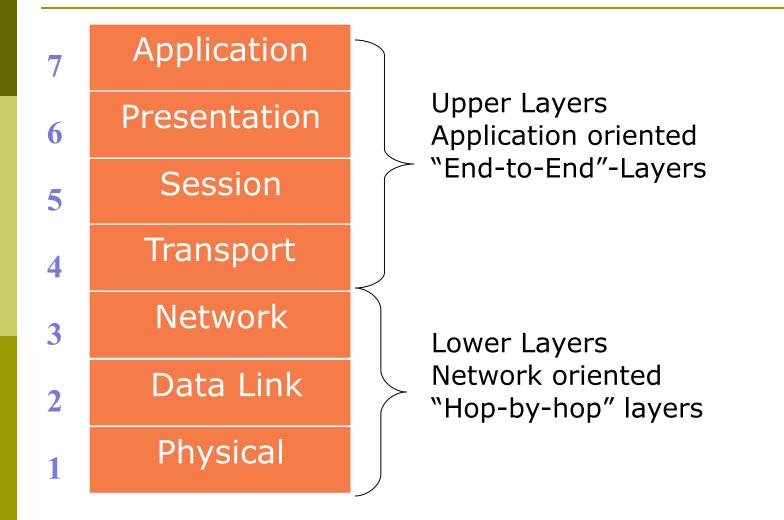
- explicit structure allows identification, relationship of complex system's pieces
  - layered reference model for discussion
- Modularisation eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure does not affect rest of system

## The IP Hourglass Model

#### Application layer SMTP DNS FTP Telnet Audio Video HTTP Transport layer UDP **RTP** TCP **Network** layer IP ADSL PPP ATM Optics Satellite **3**G Ethernet

#### Physical and Data link layer

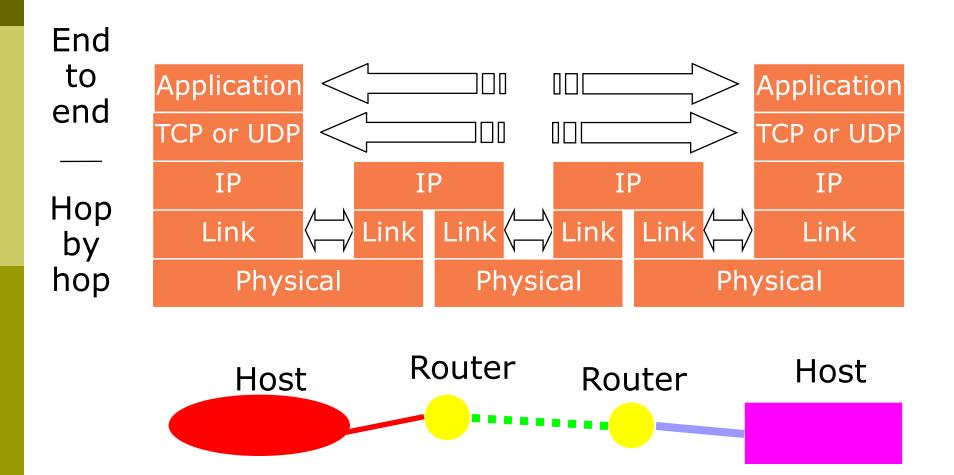
#### The OSI Model



#### OSI Model and the Internet

- Internet protocols are not directly based on the OSI model
- However, we do often use the OSI numbering system. You should at least remember these:
  - Layer 7: Application
  - Layer 4: Transport (e.g. TCP, UDP)
  - Layer 3: Network (IP)
  - Layer 2: Data link
  - Layer 1: Physical

Layer Interaction: TCP/IP Model



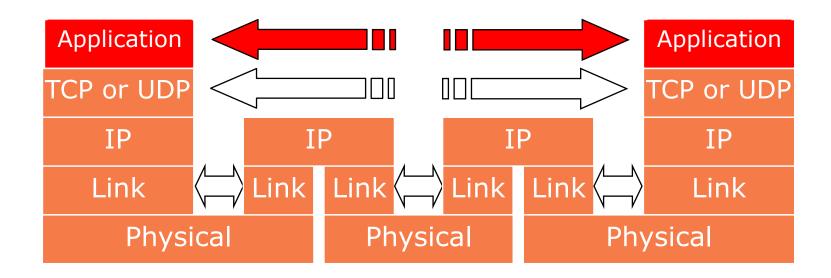
#### End-to-end layers

- Upper layers are "end-to-end"
- Applications at the two ends behave as if they can talk directly to each other
- They do not concern themselves with the details of what happens in between

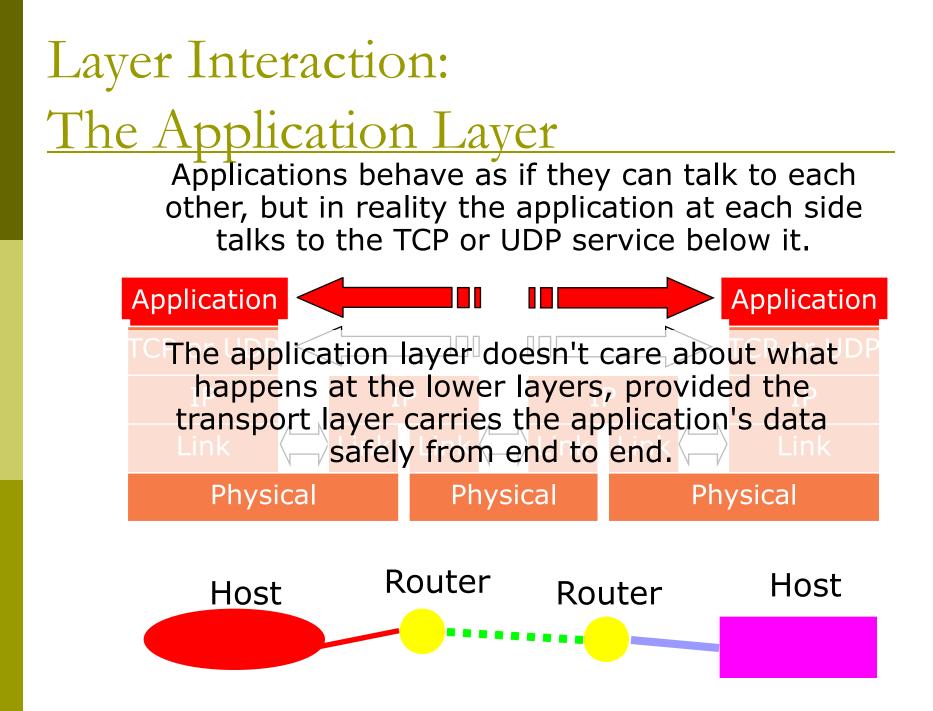
#### Hop-by-hop layers

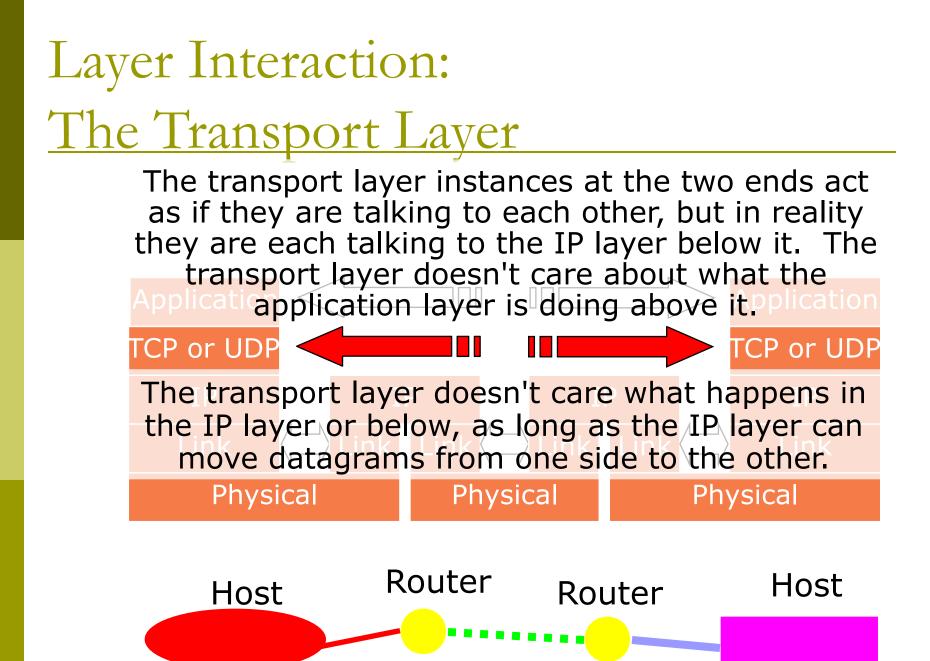
- At the lower layers, devices share access to the same physical medium
- Devices communicate directly with each other
- The network layer (IP) has some knowledge of how many small networks are interconnected to make a large internet
- Information moves one hop at a time, getting closer to the destination at each hop

Layer Interaction: TCP/IP Model









Layer Interaction: The Network Layer (IP)

The IP layer has to know a lot about the topology<br/>of the network (which host is connected to which<br/>router, which routers are connected to each<br/>other), but it doesn't care about what happens at<br/>the upper layers.IPIPIPIPIPIPThe IP layer works forwards messages hop by hop<br/>from one side to the other side.<br/>PhysicalIP sical

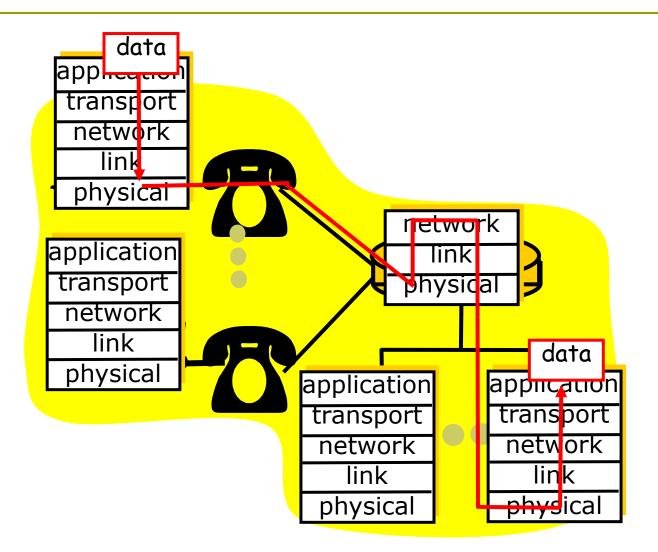


Layer Interaction: Link and Physical Layers

> The link layer doesn't care what happens above it, but it is very closely tied to the physical layer below it. All links are independent of each other, and have no way of communicating with each other. Link Link Link Link Link Link Link Physical Physical Physical



#### Layering: physical communication

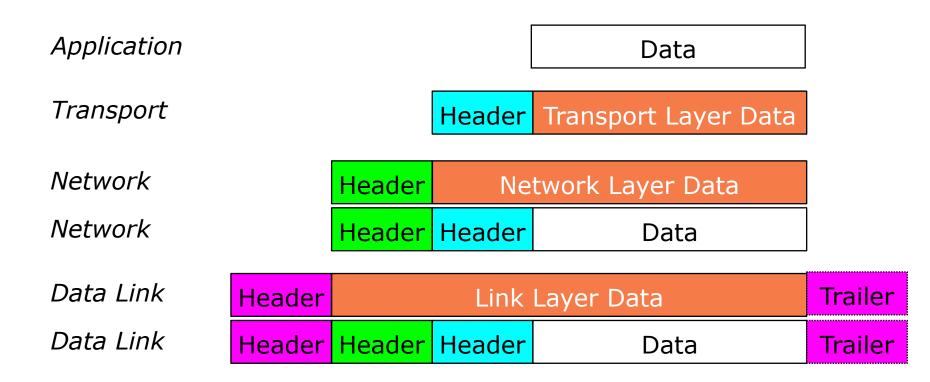


### Frame, Datagram, Segment, Packet

- Different names for packets at different layers
  - Ethernet (link layer) frame
  - IP (network layer) datagram
  - TCP (transport layer) segment
- Terminology is not strictly followed
  - we often just use the term "packet" at any layer

#### Encapsulation & Decapsulation

Lower layers add headers (and sometimes trailers) to data from higher layers



#### Layer 2 - Ethernet frame

Preamble	Dest	Source	Туре	Data	CRC
	6 bytes	6 bytes	2 bytes	46 to 1500 bytes	4 bytes

- Destination and source are 48-bit MAC addresses (e.g., 00:26:4a:18:f6:aa)
- Type 0x0800 means that the "data" portion of the Ethernet frame contains an IPv4 datagram. Type 0x0806 for ARP. Type 0x86DD for IPv6.
- Data" part of layer 2 frame contains a layer 3 datagram.

## Layer 3 - IPv4 datagram

Version IHL	Diff Services	Total Length				
Identif	ication	Flags	Frag	agment Offset		
Time to Live	Protocol	Header Checksum				
Source Address (32-bit IPv4 address)						
Destination Address (32-bit IPv4 address)						
	Padding					
Data (contains layer 4 segment)						

- Version = 4
   If no options, IHL = 5
   Source and
   Destination are 32-bit
   IPv4 addresses
- Protocol = 6 means data portion contains a TCP segment.
   Protocol = 17 means UDP.

### Layer 4 - TCP segment

Source Port				Destination Port					
	Sequence Number								
Acknowledgement Number									
Data Offset	Reserved	U R G	A C K	E O L	R S T	S Y N	F I N	Win	dow
Checksum							Urgent Pointer		
Options					าร		Padding		
Data (contains application data)									

- Source and Destination are 16-bit TCP port numbers (IP addresses are implied by the IP header)
- If no options, Data Offset = 5 (which means 20 octets)



